An Introduction to Formal Argumentation Theory



PRIFYSGOL

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Research Agenda: Argumentation for Explainable Inference

- Formal Argumentation has its roots in Non Monotonic Reasoning (NMR)
- The idea is how to reason with "rules of thumb"
- Many formalisms for NMR can be captured by formal argumentation
- One of the aims of formal argumentation is to bridge the gap between human reasoning and automated inference

An Example [Prakken]

- Paul: My car is very safe.
- Olga: Why?
- Paul: Since it has an airbag.
- Olga: It is true that your car has an airbag, but I do not think that this makes your car safe, because airbags are unreliable: the newspapers had several reports on cases where airbags did not work.
- Paul: I also read that report but a recent scientific study showed that cars with airbags are safer than cars without airbags, and scientific studies are more important than newspaper reports.

Arguments and attacks

Argument: expresses one or more reasons that lead to a proposition $a, b, c \Rightarrow d$ or $a, b \Rightarrow c; c \Rightarrow d$

An argument can *attack* another argument rebutting attack: attack one of the *conclusions* of the other argument: e, f, g $\Rightarrow \neg d$ against a, b, c $\Rightarrow d$ undercutting attack: attack the *reasons* of the other argument e, f, g \Rightarrow [a, b, c \Rightarrow d] against a, b, c $\Rightarrow d$

A: My car is very safe, since it has an airbag: has_airbag \Rightarrow safe

B: The newspapers say that airbags are not reliable, so having an airbag is not a good reason why your car is safe say(npr, ¬rel(airbag)) ⇒ ¬rel(airbag) ¬rel(airbag) ⇒ [has_airbag ⇒ safe]

C: Scientific reports say that airbags are reliable. say(sr, rel(airbag)) \Rightarrow rel(airbag)



















Next Topic

How do we evaluate an argumentation framework?

("argumentation semantics")

Argument Labellings

Each argument is labelled in, out or undec

an argument is in \Leftrightarrow all its attackers are out

an argument is **out** \Leftrightarrow it has an attacker that is in

an argument is undec ⇔ not all its attackers are out and it does not have an attacker that is in


































































ABCDE ABCDE ABCDE

"maximal": there is no other that has the same plus something "minimal": there is no other that has the same minus something



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Complete, Stable, Preferred, Grounded and Semi-Stable Labellings

in ⇔ all attackers are out out ⇔ there is an attacker that is in undec ⇔ not all attackers are out, and no attacker is in

restriction on compl. labeling no further restrictions empty undec maximal in maximal out maximal undec minimal in minimal out minimal undec

Dung-style semantics

complete semantics stable semantics preferred semantics preferred semantics grounded semantics grounded semantics semi-stable semantics

Complete Labellings and Admissible Labellings

complete labelling: in ⇔ all attackers are out out ⇔ there is an attacker that is in undec ⇔ not all attackers are out, and no attacker is in

Complete Labellings and Admissible Labellings

complete labelling: in ⇒ all attackers are out out ⇒ there is an attacker that is in undec ⇒ not all attackers are out, and no attacker is in

Complete Labellings and Admissible Labellings

admissible labelling: in ⇒ all attackers are out out ⇒ there is an attacker that is in

Roundup Labelling-Based Argumentation Semantics

admissible labelling: in ⇒ all attackers are out out ⇒ there is an attacker that is in

complete labelling: in ⇒/⇔ all attackers are out out ⇒/⇔ there is an attacker that is in undec ⇒/⇔ not all attackers are out, and no attacker is in

grounded lab.: complete with min in / min out / max undec preferred lab.: complete with max in / max out semi-stable lab.:complete with min undec stable lab.: complete with no undec

Extension-Based Argumentation Semantics (1/2)

- Args is <u>conflict-free</u> iff
 Args does not contain A,B such that A attacks B
- Args <u>defends</u> an argument C iff for each argument B that attacks C, Args contains an argument (A) that attacks B



Extension-Based Argumentation Semantics (1/2)

- Args is <u>conflict-free</u> iff
 ∄ A,B ∈ Args: A attacks B
- Args <u>defends</u> an argument C iff
 ∀ B that attacks C:
 ∃ A∈Args: A attacks B
- F(Args) = all arguments defended by Args
- Args⁺ = { A | $\exists B \in Args: B \text{ attacks } A$ }





What are $\{B\}^+$ and $F(\{B\})$ answer: $\{A,C\}$ and $\{B,D\}$

Extension-Based Argumentation Semantics (2/2)

- A set of arguments Args is called:
- <u>admissible</u> iff Args is conflict-free and Args \subseteq F(Args)
- a <u>complete extension</u> iff
 Args is conflict-free and Args = F(Args)
- a grounded extension iff
 Args is the minimal complete extension
- a <u>preferred extension</u> iff
 Args is a maximal admissible set
- a <u>stable extension</u> iff Args is a conflict-free set that attacks everything not in it
- a semi-stable extension iff Args is an admissible set with Args \cup Args⁺ maximal

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Complete	Grounded	Preferred	Stable	Semi-Stable
Extensions:	Extension:	Extensions:	Extensions:	Extensions:
Ø	Ø			
{A}		{A}		
{B,D}		{B,D}	{B,D}	{B,D}

Extension-Based semantics vs. Labelling-Based Semantics

Let Lab = (in(Lab), out(Lab), undec(Lab)) Let Args be a conflict-free set of arguments.

We define:

- Lab2Ext(Lab) = in(Lab)
- Ext2Lab(Args) = (Args, Args⁺, Ar \ (Args \cup Args⁺))

It holds that:

- If Lab is a complete labelling, then Lab2Ext(Lab) is a complete extension
- If Args is a complete extension, then Ext2Lab(Args) is a complete labelling
- For complete labellings/extensions
 Lab2Ext and Ext2Lab are each other's inverse functions

Extension-Based semantics vs. **Labelling-Based Semantics**

complete labelling \equiv complete extension grounded labelling \equiv grounded extension preferred labelling \equiv preferred extension <u>semi-stable labelling</u> \equiv semi-stable extension stable labelling \equiv stable extension

(equivalence through Lab2Ext and Ext2Lab)

<u>take home message:</u> An extension is the in-labelled part of a labelling

Next Topic

What is the relation between argumentation semantics and discussion?

Grounded Discussion Game Preferred Discussion Game

Next Topic

The Grounded Discussion Game

two players:

- proponent (P)
- opponent (O)

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- proponent (P)
- opponent (O)
- four moves $(A, B \in Ar)$ (1) P: HTB(A) "A has to be the case" (2) O: CB(B) "B can be the case" (3) O: CONCEDE(A) "I now agree that A" (4) O: RETRACT(B)
 - "I no longer hold that B can be the case"

- two players:
- proponent (P)
- opponent (O)
- four moves $(A, B \in Ar)$
- (1) P: HTB(A)
 - "A is labelled in by every complete"
- (2) O: CB(B)
 - "B can be the case"
- (3) O: CONCEDE(A)
 - "I now agree that A"
- (4) O: RETRACT(B)
 - "I no longer hold that B can be the case"

- two players:
- proponent (P)
- opponent (O)
- four moves (A,B \in Ar)
- (1) P: HTB(A)
 - "A is labelled in by every complete"
- (2) O: CB(B)
 - "Perhaps B isn't labelled out by every complete"
- (3) O: CONCEDE(A)
 - "I now agree that A"
- (4) O: RETRACT(B)
 - "I no longer hold that B can be the case"

- two players:
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- four moves (A,B \in Ar)
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 - "I now agree that A is labelled in by every complete"
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"Perhaps B isn't labelled out by every complete"

- (3) O: CONCEDE(A)
 - "I now agree that A is labelled in by every complete"
- (4) O: RETRACT(B)

"I now think that B is labelled out by every complete"





P: HTB(C)



P: HTB(C) O: CB(B)



P: HTB(C) O: CB(B) P: HTB(A)



P: HTB(C) O: CB(B) P: HTB(A) O: CONCEDE(A)



P: HTB(C) O: CB(B) P: HTB(A) O: CONCEDE(A) O: RETRACT(B)



P: HTB(C) O: CB(B) P: HTB(A) O: CONCEDE(A) O: RETRACT(B) O: CONCEDE(C)

- preceding move was CB(B) and A attacks B
- no CONCEDE or RETRACT move is applicable

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- preceding move was not a CB statement
- B has not yet been RETRACTed
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- ...of which every attacker has been RETRACTed
- and CONCEDE(A) hasn't yet been moved

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- ...of which every attacker has been RETRACTed
- and CONCEDE(A) hasn't yet been moved
- RETRACT(B)
- there has been a CB(B) statement in the past...
- ...of which an attacker has been CONCEDEd
- and RETRACT(B) hasn't yet been moved

HTB(A) This is either the first move, or

- preceding move was CB(B) and A attacks B
- no CONCEDE or RETRACT move is applicable CB(B)
- B attacks the last HTB(A) statement that is not yet CONCEDEd
- preceding move was not a CB statement
- B has not yet been RETRACTed
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- there has been a HTB(A) statement in the past...
- ...of which every attacker has been RETRACTed
- and CONCEDE(A) hasn't yet been moved RETRACT(B)
- there has been a CB(B) statement in the past...
- ...of which an attacker has been CONCEDEd
- and RETRACT(B) hasn't yet been moved

additional condition for all moves no HTB-CB repeats have occurred









P: HTB(D)





P: HTB(D) O: CB(C)





P: HTB(D) O: CB(C) P: HTB(B)

P: HTB(D) O: CB(C) P: HTB(B) O: CB(A)









P: HTB(D) O: CB(C) P: HTB(B) O: CB(A)

Proponent cannot move anymore, so discussion is terminated. Main claim not CONCEDEd, so proponent loses.









P: HTB(A) O: CB(B)



P: HTB(A) O: CB(B) P: HTB(C)

P: HTB(A) O: CB(B) P: HTB(C) O: CB(D)





P: HTB(A) O: CB(B) P: HTB(C) O: CB(D) P: HTB(A)







P: HTB(A)
O: CB(B)
P: HTB(C)
O: CB(D)
P: HTB(A)

Now we have a HTB(A)-HTB(A) repeat, so all following moves are blocked: discussion terminated.

> Main claim not CONCEDEd, so proponent loses

- A discussion is *terminated* iff no next move is possible
- A terminated discussion (starting with HTB(A)) is won by the proponent iff the opponent has moved CONCEDE(A)
- <u>soundness:</u> if a discussion is won by the proponent, then the main argument is in in the grounded labelling
- <u>completeness</u>: if an argument is in the grounded labelling then the proponent has a winning strategy for it















P: HTB(C) O: CB(B)





P: HTB(C) O: CB(B) P: HTB(E)

P: HTB(C) O: CB(B) P: HTB(E) O: CB(D)





P: HTB(C) O: CB(B) P: HTB(E) O: CB(D) P: HTB(C)




P: HTB(C) O: CB(B) P: HTB(E) O: CB(D)









P: HTB(C) O: CB(B) P: HTB(E)





P: HTB(C) O: CB(B)





P: HTB(C) O: CB(B) P: HTB(A)

P: HTB(C) O: CB(B) P: HTB(A) O: CONCEDE(A)





P: HTB(C) O: CB(B) P: HTB(A) O: CONCEDE(A) O: RETRACT(B)





P: HTB(C) O: CB(B) P: HTB(A) O: CONCEDE(A) O: RETRACT(B) O: CONCEDE(C)







- Paul: My car is very safe, because it has an airbag.
- Olga: I do not think having an airbag makes your car safe, because airbags are unreliable: the newspapers had several reports on cases where airbags did not work.
- Paul: A recent scientific study shows that airbags are in fact 99% reliable, and scientific studies are more important than newspaper reports.
- Olga: OK, I admit that airbags are reliable. However, your car is not very safe, since its maximum speed is much too high.

- P: My car is very safe, because it has an airbag.
- O: I do not think having an airbag makes your car safe, because airbags are unreliable: the newspapers had several reports on cases where airbags did not work.
- P: A recent scientific study shows that airbags are in fact 99% reliable, and scientific studies are more important than newspaper reports.
- O: OK, I admit that airbags are reliable. However, your car is not very safe, since its maximum speed is much too high.

- P: HTB(airbag \Rightarrow safe)
- O: I do not think having an airbag makes your car safe, because airbags are unreliable: the newspapers had several reports on cases where airbags did not work.
- P: A recent scientific study shows that airbags are in fact 99% reliable, and scientific studies are more important than newspaper reports.
- O: OK, I admit that airbags are reliable. However, your car is not very safe, since its maximum speed is much too high.

- P: HTB(airbag \Rightarrow safe)
- O: $CB(npr \Rightarrow \neg airbagrel \Rightarrow \neg [airbag \Rightarrow safe])$
- P: A recent scientific study shows that airbags are in fact 99% reliable, and scientific studies are more important than newspaper reports.
- O: OK, I admit that airbags are reliable. However, your car is not very safe, since its maximum speed is much too high.

- P: HTB(airbag \Rightarrow safe)
- O: $CB(npr \Rightarrow \neg airbagrel \Rightarrow \neg [airbag \Rightarrow safe])$
- P: HTB(study \Rightarrow airbagrel)
- O: OK, I admit that airbags are reliable. However, your car is not very safe, since its maximum speed is much too high.

- P: HTB(airbag \Rightarrow safe)
- O: $CB(npr \Rightarrow \neg airbagrel \Rightarrow \neg [airbag \Rightarrow safe])$
- P: HTB(study \Rightarrow airbagrel)
- O: CONCEDE(study \Rightarrow airbagrel) However, your car is not very safe, since its maximum speed is much too high.

- P: HTB(airbag \Rightarrow safe)
- O: $CB(npr \Rightarrow \neg airbagrel \Rightarrow \neg [airbag \Rightarrow safe])$
- P: HTB(study \Rightarrow airbagrel)
- O: CONCEDE(study \Rightarrow airbagrel) RETRACT(npr $\Rightarrow \neg$ airbagrel $\Rightarrow \neg$ [airbag \Rightarrow safe])

- P: HTB(airbag \Rightarrow safe)
- O: $CB(npr \Rightarrow \neg airbagrel \Rightarrow \neg [airbag \Rightarrow safe])$
- P: HTB(study \Rightarrow airbagrel)
- O: CONCEDE(study \Rightarrow airbagrel) RETRACT(npr $\Rightarrow \neg$ airbagrel $\Rightarrow \neg$ [airbag \Rightarrow safe]) CB(highspeed $\Rightarrow \neg$ safe)

Next Topic

The Preferred Discussion Game

Socratic Discussion

Answer me this. As soon as one man loves another, which of the two becomes the friend? the lover of the loved, or the loved of the lover? Or does it make no difference?

None in the world, that I can see

How? Are both friends, if only one loves?

I think so

- *Indeed!* is it not possible for one who loves, not to be loved in return (...) ? It is.
- *Nay, is it not possible for him even to be hated? (...) Don't you believe this to be true?* Quite true.

Well, in such a case as this, the one loves, the other is loved. Just so.

Which of the two, then, is the friend of the other? The lover of the loved, whether or not he be loved in return, and even if he be hated, or the loved of the lover? or is neither the friend of he other, unless both love each other?

The latter certainly seems to be the case, Socrates.

If so, I continued, we think differently now from what we did before. (...) Yes, I'm afraid we have contradicted ourselves.

Traditional Dialogue vs. Socratic Dialogue

P: claim tr *"I think that there will be a tax relief."*O: why tr *"Why do you think so?"*P: because pmp ⇒ tr *"Because of the fact that the politicians made a promise."*O: concede tr *"OK, you are right."*

Traditional Dialogue vs. Socratic Dialogue

- P: claim tr
 - *"I think that tr."*
- O: but-then tr \Rightarrow bd *"Then you implicitly also hold that bd."*
- P: concede bd "Yes I do."
- O: but-then bd \Rightarrow feu "Then you implicitly also hold that feu."
- P: concede feu "Yes I do."
- O: but-then feu $\Rightarrow \neg$ tr

"Then you implicitly also hold that ¬tr."

P: concede ¬tr "Oops, you're right; I caught myself in..."

"because" versus "but-then"



reasoning goes backward

proponent constructs path

originates from true

both parties become committed

reasoning goes forward opponent constructs path leads to *false* only proponent

becomes committed

<u>definition</u> admissible labelling: if argument is in then all its attackers are **out** if argument is **out** then it has an attacker that is in

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proposition

An argument is in a preferred extension iff it is in a complete extension iff it is in an admissible set iff it is labelled in by an admissible labelling





M: in(D) "I have an admissible labelling in which D is labelled in."



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(1) Each move of M (except the first) contains an attacker of the directly preceding move of S.

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S: out(C) "But then in your labelling it must also be the case that D's attacker C is labelled out. Based on which grounds?"
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S: out(A) "But then in your labelling it must also be the case that B's attacker A is labelled out. Based on which grounds?"
M: in(B) "A is labelled out because B is labelled in."

(2) Each move of S contains an attacker of <u>some</u> previous move of M.

M: in(D) "I have an admissible labelling in which D is labelled in."
S: out(C) "But then in your labelling it must also be the case that D's attacker C is labelled out. Based on which grounds?"
M: in(B) "C is labelled out because B is labelled in."
S: out(A) "But then in your labelling it must also be the case that B's attacker A is labelled out. Based on which grounds?"
M: in(B) "A is labelled out because B is labelled in."

(3) S is not allowed to repeat his moves.

M: in(D) "I have an admissible labelling in which D is labelled in."
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S: out(A) "But then in your labelling it must also be the case that B's attacker A is labelled out. Based on which grounds?"
M: in(B) "A is labelled out because B is labelled in."

(4) M is allowed to repeat his moves.




M: in(E) "I have an admissible labelling in which E is labelled in."



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M: in(C) "D is labelled out because C is labelled in."



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M: in(C) "D is labelled out because C is labelled in."
S: out(E) "But then in your labelling it must also be the case that C's attacker E is labelled out. This <u>contradicts</u> with your earlier claim that E is labelled in."

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(5) If S uses an argument previously used by M, then S wins the discussion.

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S: out(D) "But then in your labelling it must also be the case that E's attacker D is labelled out. Based on which grounds?"
M: in(C) "D is labelled out because C is labelled in."
S: out(E) "But then in your labelling it must also be the case that C's attacker E is labelled out. This <u>contradicts</u> with your earlier claim that E is labelled in."

(6) If M uses an argument previously used by S, then S wins the discussion.

M: in(E) "I have an admissible labelling in which E is labelled in."
S: out(D) "But then in your labelling it must also be the case that E's attacker D is labelled out. Based on which grounds?"
M: in(C) "D is labelled out because C is labelled in."
S: out(E) "But then in your labelling it must also be the case that C's attacker E is labelled out. This <u>contradicts</u> with your earlier claim that E is labelled in."

(7) If M cannot make a move anymore, then S wins the discussion.

M: in(E) "I have an admissible labelling in which E is labelled in."
S: out(D) "But then in your labelling it must also be the case that E's attacker D is labelled out. Based on which grounds?"
M: in(C) "D is labelled out because C is labelled in."
S: out(E) "But then in your labelling it must also be the case that C's attacker E is labelled out. This <u>contradicts</u> with your earlier claim that E is labelled in."

(8) If S cannot make a move anymore, then M wins the discussion.

<u>THEOREM</u>

Argument A is labelled in by at least one admissible labelling iff M can win the Socratic discussion game (for A).

<u>THEOREM</u>

Argument A is in at least one preferred extension iff M can win the Socratic discussion game (for A).

Complete Semantics as Socratic Discussion

<u>THEOREM</u>

Argument A is in at least one complete extension iff M can win the Socratic discussion game (for A).

Skeptical Complete vs Credulous Complete

- Skeptical Complete (= grounded) The argument is accepted in every reasonable position (complete labelling) "Therefore you have to hold that..." persuasion dialogue
- Credulous Complete (= credulous preferred) The argument is accepted in at least one reasonable position (complete labelling) "Therefore I can hold that..."
 <u>Socratic dialogue</u>

• classical logic:

• argumentation:

- classical logic: based on notion of <u>truth</u> (entails what is model-theoretically true)
- argumentation:

- classical logic: based on notion of <u>truth</u> (entails what is model-theoretically true)
- argumentation: based on notion of *justification* (entails what can be defended in rational discussion)

- classical logic: based on notion of <u>truth</u> (entails what is model-theoretically true)
- argumentation: based on notion of *justification* (entails what can be defended in rational discussion)
- discussions can be used by the system to <u>explain</u> its answer to the user
- different semantics express different <u>types</u> of rational discussion
- allows (in principle) for <u>dynamic</u> and <u>user-based</u> updating of the underlying knowledge base

Next Topic

Rationality Postulates

If we select arguments using only the structure of the graph, then how do we know their conclusions make sense?

knowledge base

argumentation framework (1) argument (+attack) construction knowledge base





From Arguments to Conclusions

labelling of args → labelling of concs for each conclusion, give it the label of the "best" (highest label) argument that yields it

How Things Go Wrong (1/5)

rr \Rightarrow mm \rightarrow hspp \Rightarrow bb \rightarrow \neg hs

 $A1 = (r) \Rightarrow m$ $A2 = (p) \Rightarrow b$ $A3 = A1 \rightarrow hs$ $A4 = A2 \rightarrow \neg hs$

Conclusions m and b are justified under any semantics but what about hs and ¬hs?

How Things Go Wrong (2/5)

r $r \Rightarrow m$ $m \supset hs$ p $p \Rightarrow b$ $b \supset \neg hs$

$$("\rightarrow" \equiv "\vdash")$$

A1: (r) \Rightarrow m A2: (p) \Rightarrow b A3: (A1, m \supset hs) \rightarrow hs A4: (A2, b \supset \neg hs) \rightarrow \neg hs A5: (A3, b \supset \neg hs) \rightarrow \neg b A6: (A4, m \supset hs) \rightarrow \neg m

So far, so good...

How Things Go Wrong (3/5)





How Things Go Wrong (4/5)



- Grounded semantics: no justified arguments
- Why not use preferred or stable semantics?
- Reiter and Pollock also do this...

How Things Go Wrong (5/5)

John: "Cup of coffee contains sugar." Mary: "Cup of coffee doesn't contain sugar." John: "I'm unreliable." Mary: "I'm unreliable." Weather Forecaster: "Tomorrow rain."



Rationality Postulates

Let J be a set of conclusions yielded by an argumentation formalism.

- <u>direct consistency</u>
 J does not contain contraries (p and ¬p)
- <u>closure</u>
 J is closed under the strict rules
- indirect consistency the closure of J under strict rules is directly consistent
- <u>crash-resistance</u>

no set of formulas can make a totally unrelated set of formulas completely irrelevant, when being merged to it

<u>non-interference</u>

no set of formulas can influence the entailment of a totally unrelated set of formulas, when being merged to it

Rationality Postulates

direct consistency closure indirect consistency

Caminada & Amgoud AIJ 2007

crash-resistance non-interference backwards compatibility

Caminada, Dunne & Carnielli JLC 2011

Transposition

Take the following strict rule: $a_1, ..., a_{i-1}, a_i, a_{i+1}, ..., a_n \rightarrow c$ A transposition of this rule is: $a_1, ..., a_{i-1}, \neg c, a_{i+1}, ..., a_n \rightarrow \neg a_i$ (for some $1 \le i \le n$)

A set of strict rules S is closed under transposition iff it contains all transpositions of the rules in S.

Restricted versus unrestricted rebut

$((a) \Rightarrow b) \Rightarrow c$ $((d) \Rightarrow e) \Rightarrow \neg c$

Restricted versus unrestricted rebut

$$((a) \rightarrow b) \rightarrow c$$
$$((d) \Rightarrow e) \Rightarrow \neg c$$
Restricted versus unrestricted rebut

$$((a) \Rightarrow b) \rightarrow c$$
$$((d) \rightarrow e) \Rightarrow \neg c$$

Restricted versus unrestricted rebut

$$\begin{array}{l} ((a) \Rightarrow b) \rightarrow c \\ ((d) \rightarrow e) \Rightarrow \neg c \end{array}$$

unrestricted rebut: an argument can be rebutted on a conclusion derived by at least one defeasible rule

restricted rebut:

an argument can be rebutted only on the direct consequent of a defeasible rule

Satisfying the rationality postulates (1/2)

Satisfying the rationality postulates (1/2)

When using strict rules as primitives:

Satisfying the rationality postulates (1/3)

When using strict rules as primitives:

- Step 1:
 - close the strict rules under transposition
 - use *restricted rebut* for defining the attack relation
- Step 2:
 - use a <u>complete-based</u> semantics that yields at least one extension/labelling
- Step 3: (standard)

Satisfying the rationality postulates (2/3)

When using strict rules as primitives:

- Step 1:
 - close the strict rules under transposition
 - use *unrestricted rebut* for defining the attack relation
- Step 2:
 - use grounded semantics

 Step 3: (standard)

Satisfying the rationality postulates (2/2)

When using strict rules as classical entailment:

Satisfying the rationality postulates (3/3)

When using strict rules as classical entailment:

- Step 1:
 - use restricted rebut for defining the attack relation
 - remove all inconsistent arguments
- Step 2:
 - use a <u>complete-based</u> semantics that yields at least one extension/labelling
- Step 3: (standard)

 $S = \{ w1; w2; w3; \\ D = \{ w1 \Rightarrow b1; w2 \Rightarrow b2; w3 \Rightarrow b3 \}$

 $S = \{ w1; w2; w3; b2, b3 \rightarrow \neg b1; b1, b3 \rightarrow \neg b2; b1, b2 \rightarrow \neg b3 \}$ $D = \{ w1 \Rightarrow b1; w2 \Rightarrow b2; w3 \Rightarrow b3 \}$

 $S = \{ w1; w2; w3; b2, b3 \rightarrow \neg b1; b1, b3 \rightarrow \neg b2; b1, b2 \rightarrow \neg b3 \}$ $D = \{ w1 \Rightarrow b1; w2 \Rightarrow b2; w3 \Rightarrow b3 \}$

 $A_{1}: (w1) \Rightarrow b1$ $A_{2}: (w2) \Rightarrow b2$ $A_{3}: (w3) \Rightarrow b3$

 $S = \{ w1; w2; w3; b2, b3 \rightarrow \neg b1; b1, b3 \rightarrow \neg b2; b1, b2 \rightarrow \neg b3 \}$ $D = \{ w1 \Rightarrow b1; w2 \Rightarrow b2; w3 \Rightarrow b3 \}$

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Why Admissibility

 $S = \{ w1; w2; w3; b2, b3 \rightarrow \neg b1; b1, b3 \rightarrow \neg b2; b1, b2 \rightarrow \neg b3 \}$ $D = \{ w1 \Rightarrow b1; w2 \Rightarrow b2; w3 \Rightarrow b3 \}$



Why Admissibility

 $S = \{ w1; w2; w3; b2, b3 \rightarrow \neg b1; b1, b3 \rightarrow \neg b2; b1, b2 \rightarrow \neg b3 \}$ $D = \{ w1 \Rightarrow b1; w2 \Rightarrow b2; w3 \Rightarrow b3 \}$

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Research Challenge: finding the magic combinations

extensions (labellings) of conclusions

(3) determining status of conclusions extensions (labellings) of arguments

(2) applying argumentation semantics

argumentation framework

(1) argument (+attack) construction

knowledge

base

Open Research Questions:

 How to satisfy crash resistance when strict rules are generated by classical logic and defeasible rules have different strength?
How to satisfy closure and consistency when applying unrestricted rebut and defeasible rules have different strength?

Further Reading:

M.W.A. Caminada Rationality Postulates: applying argumentation theory for non-monotonic reasoning (Handbook of Formal Argumentation)

Recent Developments: Heyninck & Straßer IJCAI17 / AAMAS19

A < B iff DefRules(A) $\neq \emptyset$ and $\forall d_B \in min(DefRules(B)) \exists d_A \in DefRules(A): d_A < d_B$

Generalized Rebut:

A GeRe B iff Conc(A) = $\neg b_1 \lor \ldots \lor \neg b_n$ (n ≥ 1) and each b_i occurs in B

Example: $\Rightarrow \neg b \lor \neg d$ rebuts (GeRe) (($\Rightarrow b$) $\Rightarrow c$) $\Rightarrow d$

A defeats B iff A GeRe B and A \prec B

This approach satisfies all rationality postulates (direct/indirect consistency, closure, crash-resistance, non-interference) but doesn't (yet) support undercutting (rebut only)

Argumentation Research: the Good, the Bad and the Ugly

<u>The Good</u>

research on how to to meaningfully draw conclusions based on (sufficiently rich) arguments and how to explain it

Argumentation Research: the Good, the Bad and the Ugly

<u>The Bad</u>

research that only provides an abstraction without specifying what would be the full theory

always ask: "What is it that your abstract theory provides an abstraction of"

Argumentation Research: the Good, the Bad and the Ugly

The Ugly

research that uses argumentation for a purpose that could be done in a much simpler way (e.g. MCS and LP)

Argumentation versus Machine Learning Only

Why not use ML on a big corpus of reasoning cases?

1) because the corpus may contain <u>flawed reasoning</u> (e.g. biases or logical errors)

2) because you'd need to have a <u>huge corpus</u>, which might need to be generated

3) because you want to have <u>explainability</u>

idea: use argument mining to obtain the argument schemes

Research Challenges

- How to satisfy closure, consistency and crash-resistance under different settings?
- Does exposure to argument-based discussion increase the user's confidence in the system's inferences?
- What would the discussion look like if the moves are <u>rules</u> instead of <u>arguments</u>?

Further Reading (on topics in this presentation)

- Argumentation semantics
 - Dung, AIJ 1995 (landmark paper)
 - Baroni, Caminada & Giacomin (HOFA 2018, Ch4)
 - Caminada & Dunne (Argument&Computation 2019)
- Argumentation discussion games
 - Caminada (HOFA 2018, Ch10)
- Instantiated argumentation & rationality postulates
 - Modgil & Prakken (HOFA 2018, Ch6)
 - Caminada (HOFA 2018, Ch15)
 - Heyninck & Straßer (IJCAI 2017 / AAMAS 2019)
- Warnings against pure abstract argumentation
 - Caminada & Wu (BNAIC 2011)
 - Prakken & de Winter (COMMA 2018)

Further Reading (on topics <u>not</u> in this presentation)

- How argumentation captures other forms of NMR:
 - Default Logic: Dung (AIJ 1995), Caminada et al (JLC 2012)
 - Logic Programming: Dung (AIJ 1995), Caminada & Schulz (JAIR 2017)
 - Classical Logic and Maximal Consistent Sets: Besnard & Hunter (HOFA 2018 Ch9)
- Argument Schemes:
 - Macagno et al (HOFA 2018 Ch11)
- Formal argumentation and human intuitions:
 - Yu, Xu & Liao (Studies in Logic, 2018)
 - Cramer & Guillaume (COMMA 2018 / JELIA 2019)

Further Reading (on topics <u>not</u> in this presentation)

- Argument Mining
 - Cabrio & Villata (IJCAI 2018)
 - IJCAI 2019 tutorial Federico Cerutti (tomorrow)
- Argumentation implementations
 - Cerutti et al (HOFA 2018, Ch14)
 - Nofal et al (AIJ 2014)
 - ICCMA 2015/2017/2019
- Online demonstrators:
 - DISCO: http://disco.cs.cf.ac.uk/
 - TOAST: http://toast.arg-tech.org/
Handbook of Formal Argumentation



Editors: Baroni, Gabbay, Giacomin & van der Torre College Publications (2018) 1028 pages \$33.75 / €30.16 / £27.55

More on Argumentation at IJCAI

- Tutorial Federico Cerutti *Argumentation and Machine Learning* Monday 8.30am – 12.30pm
- Technical session
 Computational Models of Argument Friday 9.30am 10.30am