

Extending SWRL to Express Fully-Quantified Constraints

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<http://www.csd.abdn.ac.uk/research/akt/cif>



In this talk...

- We argue for the utility of an expressive quantified constraint language for the SW logic layer
 - based on classical range-restricted FOL
- We develop a quantified constraint representation as an extension of the Semantic Web Rule Language (SWRL)
 - compatible with OWL as well as RDFS
- We illustrate the use of the CIF/SWRL representation in the context of a practical SW reasoning application
 - based on the CS AKTive Space demonstrator



Context

- An approach to knowledge fusion in open, distributed environments
- Gather relevant data from multiple network sources, along with constraints on how the data can be used
- Fuse data and constraints into a dynamically-composed constraint satisfaction problem (CSP)
- Solutions (if any) are relayed back to the query originator
- Examples:
 - e-commerce: configuring custom packages
 - e-science: coalition operations for virtual organisations



Semantic Web context

- Fits W3C vision of a task-oriented Web “better enabling people and computers to work in cooperation”
- RDF(S) fits our minimal requirements for data to be expressed against a semantic data model
- OWL allows far richer modelling, and also supports ontology mapping
- SWRL it is not sufficiently expressive for our needs, and we therefore propose an extension that allows the representation of fully-quantified constraints



Application: AKTive Workgroup Builder

- Builds on CS AKTive Space demonstrator (winner of the 2003 Semantic Web Challenge)
- CAS is a large-scale repository of semantic metadata on computing science activities in the UK
- AWB uses constraints to select individuals from CAS to form working groups that satisfy particular requirements
- Could be used, for example, to form “expert panels”, suggest partners for collaborative projects, or organise workshops



CS AKTive Space

AKT CS AKTive Space
Take a tour through CS AKTive Space

[About this page](#)

 research area/region
 region/research area

Research area

- Theory of Computation
 - mathematical logic and formal languages
 - logics and meanings of programs
 - analysis of algorithms and problem compl
 - computation by abstract devices
 - general
- Mathematics of Computing
 - probability and statistics
 - discrete mathematics
 - numerical analysis
 - general
- Information Systems
 - information interfaces and presentation
 - information systems applications
 - information storage and retrieval
 - database management
 - general
- Computing Methodologies
 - document and text processing
 - simulation and modeling

Radial: 50 miles
Map: uk-political

Researcher

Top 5
 10
 20
 unlimited

Order by
 Grant total
 RAE result

- KN Brown
- P Edwards
- AD Preece
- TJF Norman
- JRW Hunter
- PMD Gray
- DH Sleeman
- EB Reiter
- MW Freeston
- PJF Lucas
- GJL Kemp



Kinds of rule

- **Derivation rules**

```
( $\forall$ ?x,?p,?s,?g) hasParent(?x,?p)  $\wedge$  hasSibling(?p,?s)  $\wedge$   
hasSex(?s,?g)  $\wedge$  (?g='male')  $\Rightarrow$  hasUncle(?x,?s)
```

- **Rewrite rules**

```
( $\forall$ ?x,?y) akt:supervises(?x,?y)  $\Rightarrow$  foaf:knows(?x,?y)
```

- **Event-condition-action rules**

```
ON REGISTER-STUDENT(?s)  
WHERE supervises(?t,?s)  $\wedge$  hasGroup(?t,?g)  
DO ASSERT(hasGroup(?s,?g))
```

- **Quantified constraints**

```
( $\forall$ ?t,?s,?g) Tutor(?t)  $\wedge$  hasStatus(?t,'research')  $\wedge$   
supervises(?t,?s)  $\wedge$  hasSubjectGrade(?s,'Computing',?g)  $\Rightarrow$   
(?g>60)
```



Constraint Interchange Format (CIF)

- A CIF constraint consists of some universally quantified implications, followed by a conjunction of predicates, possibly existentially quantified

```
( $\forall$ ?t) Tutor(?t)  $\wedge$  hasStatus(?t,'research')  $\Rightarrow$   
( $\exists$ ?s,?g) supervises(?t,?s)  $\wedge$   
hasSubjectGrade(?s,'Computing',?g)  $\wedge$   
(?g>60)
```

- Why have both types of quantifier? "Sometimes readability is more important than parsimony"
[Artificial Intelligence; A Modern Approach, Russell & Norvig, 1995]
- Choice of how to implement the constraints is left to local reasoners...



Extending SWRL to CIF/SWRL

- Incorporates SWRL constructs where possible (striving to simplify the original CIF syntax)
- Constraints are defined as quantified implications:
 - re-use the implication structure from SWRL
 - allow for nested quantified implications within the consequent
 - innermost-nested implication will have an empty body as it is always of the form *true* \Rightarrow ...
- Example in the informal, human-readable syntax:

$$\begin{aligned} &(\forall ?x \in X, ?y \in Y) p(?x, ?y) \wedge Q(?x) \Rightarrow \\ &(\forall ?z \in Z) q(?x, ?z) \wedge R(?z) \Rightarrow \\ &(\exists ?v \in V) s(?y, ?v) \end{aligned}$$


Abstract syntax (extended from SWRL)

constraint ::= 'Implies(' [*URIreference*] { *annotation* }
 quantifiers antecedent consequent)'

antecedent ::= 'Antecedent(' { *atom* })'

consequent ::= 'Consequent(' *constraint* | { *atom* })'

quantifiers ::= 'Quantifiers(' { *q-atom* })'

q-atom ::= *quantifier* '(' *q-var q-set*)'

quantifier ::= 'forall' | 'exists'

q-var ::= *i-variable*

q-set ::= *classID*



Example in the abstract syntax

$$\begin{aligned} &(\forall ?x \in X, ?y \in Y) p(?x, ?y) \wedge Q(?x) \Rightarrow \\ &(\forall ?z \in Z) q(?x, ?z) \wedge R(?z) \Rightarrow \\ &(\exists ?v \in V) s(?y, ?v) \end{aligned}$$

```
Implies (
  Quantifiers (forall (I-variable (x) X) forall (I-variable (y) Y))
  Antecedent (p (I-variable (x) I-variable (y)) Q (I-variable (x)))
  Consequent (
    Implies (
      Quantifiers (forall (I-variable (z) Z))
      Antecedent (
        q (I-variable (x) I-variable (z)) R (I-variable (z)))
      Consequent (
        Implies (
          Quantifiers (exists (I-variable (v) V))
          Antecedent ()
          Consequent (s (I-variable (y) I-variable (v)))
        )
      )
    )
  )
)
```



Sketch of CIF/SWRL RDF syntax

- New class, **cif:Constraint**; two attached properties:
 - **cif:hasQuantifiers** (range rdf:List)
 - **cif:hasImplication** (range ruleml:Imp)
- Parent class **cif:Quantifier**; two sub-classes:
 - **cif:Forall**
 - **cif:Exists**
- Two properties attached to **cif:Quantifier**:
 - **cif:var** (range rdf:Resource - URIref of SWRL variable)
 - **cif:set** (range rdf:Resource - URIref of OWL/RDFS classID)
- Note: SWRL RDF syntax allows ruleml:body to be any RDF list, so allows the nested inclusion of a **cif:Constraint**.



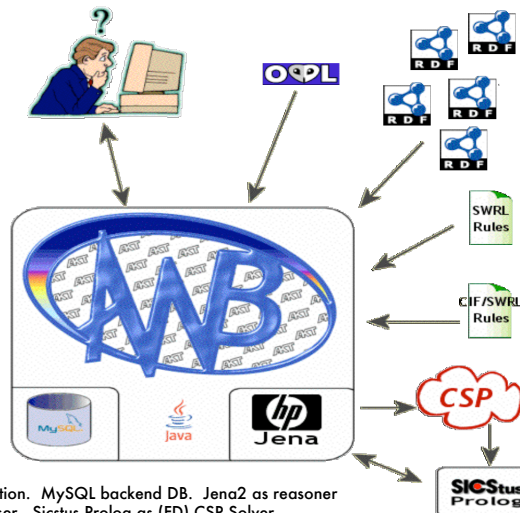
AKTive Workgroup Builder

Constructing a workgroup involves several steps:

- Defining constraints about the nature of the workgroup:
 - "generic" workgroup constraints
 - context/user-specific constraints
- Gathering (RDF) data about candidate workgroup members
- Generating entailments: ontological & from derivation rules
- Composing a CSP from the data & constraints
- Solving the CSP



AWB schematic



J2EE Application. MySQL backend DB. Jena2 as reasoner and RDF parser. Sicstus Prolog as (FD) CSP Solver.



Reasoning isn't the (only) hard part

- AWB v1 uses a cut-down version of the CS AKTive Space from ISWC 2003:
 - OWL Full **AKT Portal Ontology** was cut to OWL Lite (lack of reasoners)
 - 10M triples were reduced to a subset (querying limitations)
 - data was cleaned-up (errors, duplication, incompleteness)
- AWB v1 is aimed at scheduling AKT meetings:
 - data covers 5 AKT partners
 - reduced APO ignores OWL DL/Full constructs, flattens hierarchy
 - some (minor) fixes made to ontology and data



AWB: reasoning using SWRL

- Initially, on loading candidate instances, the AWB computes OWL Lite entailments (every **Professor-in-Academia** is a **Person-in-Academia**, etc)
- Then SWRL derivation rules compute additional facts
 - Example: "if a person has an affiliation with an organisation, and that organisation has a postal address with a city then this implies that the person has a base location of the same city"

```
( $\forall ?p, ?u, ?a, ?c$ ) Person(?p)  $\wedge$  Organisation(?u)  $\wedge$   
has-affiliation(?p, ?u)  $\wedge$  has-postal-address(?u, ?a)  $\wedge$   
address-city(?a, ?c)  $\Rightarrow$  has-base-location(?p, ?c)
```



AWB: CIF/SWRL simple example

- Example: “every workgroup must contain at least 1 member who is a Professor”:

```
( $\forall$ ?g $\in$ Workgroup)
( $\exists$ ?p $\in$ Professor-In-Academia)
  has-member(?g,?p)
```

- Notes

- this kind of existentially-quantified statement can be expressed implicitly in SWRL using OWL **someValuesFrom**
- we prefer to express all quantifiers uniformly and explicitly, and leave the reasoner the option of transforming the constraints to a suitable implementation form



In RDF...

Note: the (OWL Lite) ontology URI is represented by the entity **&akt;**

```
< cif:Constraint >
< cif:hasQuantifiers rdf:parseType="Collection" >
  < cif:Forall >
    < cif:var rdf:resource="#g" />
    < cif:set rdf:resource="&akt;#Workgroup" />
  < /cif:Forall >
  < cif:Exists >
    < cif:var rdf:resource="#p" />
    < cif:set rdf:resource="&akt;#Professor-In-Academia" />
  < /cif:Exists >
< /cif:hasQuantifiers >
< cif:hasImplication >
  < swrl:Imp >
    < swrl:body rdf:parseType="Collection" />
    < swrl:head rdf:parseType="Collection" >
      < swrl:IndividualPropertyAtom >
        < swrl:classPredicate rdf:resource="&akt;#has-member" />
        < swrl:argument1 rdf:resource="#g" />
        < swrl:argument2 rdf:resource="#p" />
      < /swrl:IndividualPropertyAtom >
    < /swrl:head >
  < /swrl:Imp >
< /cif:hasImplication >
< /cif:Constraint >
```



AWB: CIF/SWRL 2nd example

- “Any workgroup with at least 5 members must contain people from different sites”:

$$\begin{aligned} &(\forall ?g \in \text{Workgroup}) \text{has-size} (?g, ?s) \wedge (?s \geq 5) \Rightarrow \\ &(\exists ?p1, ?p2 \in \text{Person}) \text{has-member} (?g, ?p1) \wedge \\ &\quad \text{has-base-location} (?p1, ?b1) \wedge \\ &\quad \text{has-member} (?g, ?p2) \wedge \\ &\quad \text{has-base-location} (?p2, ?b2) \wedge (?b1 \neq ?b2) \end{aligned}$$

- Notes

- uses the (derived) property **has-base-location** from our SWRL example to indicate a person’s “site”
- interacts with previous SWRL derivation rule



Solving CIF

- Solving CIF constraints can be implemented by dynamically composing the constraints and available data instances into a CSP, code-generated for use with a particular finite domain solver
- Solvers used to date include
 - ECLiPSe (<http://www.icparc.ic.ac.uk/eclipse/>)
 - Sicstus Prolog FD library (<http://www.sics.se/isl/sicstus/>)
- 3 steps
 - form variable domains from candidate instance data
 - post constraints (translate CIF to native solver code)
 - label variables (instantiate vars such that constraints are satisfied...)



CIF's closed world assumption

- We are making a closed world assumption, at the time the finite domain CSP is composed
- This might seem contradictory to the general vision of an open world Semantic Web (and OWL DL)
- In practice, a finite number of candidate instances are always available at run-time, whether
 - gathered from a local cache (as in the current AWB)
 - acquired through some wider search (always "best-effort" on the Web)
- (We are essentially doing A-Box reasoning...)



CIF & OWL (DL)

- Our approach is not incompatible with the use of other reasoning mechanisms
- Example: OWL DL class restrictions can usefully be employed in CIF expressions to specify the domains of variables,
 - in the quantifier expressions (as the value of a **cif:set** property)
 - within the heads and bodies of the implications (unary-predicate **atoms**)
- We have yet to explore the computational complexities arising from this :-)



CIF & RDFS

- It is perfectly feasible to use CIF with only RDFS data models
- (This is true of SWRL as well, although of course SWRL has no way to handle existential quantification without OWL DL constructs)
- RDFS is relatively widely used on the current Semantic Web (Dublin Core, RSS, vCards, and FOAF are among the most widely-instantiated SW schemas)
- We feel this makes CIF immediately useful for practical SW applications



Statements about constraints

- The **URIreference** and **annotation** features from OWL allow statements to be made about constraints
- This allows provenance information to be attached
- It also allows other kinds of metadata specific to the usage of constraints (example: “strength” - hard/soft?)
- We use constraint reification in the solving process, where it becomes useful to reason about which constraints are currently satisfied
- Negotiation and argumentation can be used to soften (or in some cases harden) constraints



Conclusion

- **Contributions:**
 - a representation for fully-quantified constraints at the Semantic Web logic layer, as an extension to SWRL
 - a realistic test-bed application: the AKTive Workgroup Builder
- **Work on multi-strategy reasoning in the AWB is ongoing:**
 - Jena 2 for OWL Lite reasoning
 - trials with Jena 2, Hoolet, Jess, & Prolog for SWRL inference
 - calling SICStus Prolog FD library from Java via PrologBeans
- **We aim to combine these into a practical hybrid reasoner, exploring complexity/scalability trade-offs**

